ECE401/8001 VISUAL SIGNAL PROCESSING AND COMMUNICATIONS

Winter 2005 Prof.: Dr. Zhihai (Henry) He

Project Phase I: Huffman Coding 02/03/2005

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Purpose

Design a Huffman code with Matlab or C.

Procedure

Step 1: Write a function to find the probability of English characters a-z, A-Z, etc... in the article "state.txt"

Step 2: Construct the Huffman code word table.

Step 3: Compute the average length

Step 1: (approach of the Problem)

First of all I have to read the "state.txt" file into Matlab. (Using *fopen* and *fscanf*). The whole training text is stored in a string variable I named *sffile*. The number of characters is stored in a variable named *N*. In this particular file there are **45592** symbols.

Next, we have to count the symbols, and what is their frequency in the training text to find the probability distribution.

Before counting the symbols, I have to create an alphabet. I created an alphabet inside the function *enumsymbols()* where I stored all characters in a string vector. I included the "special" characters inside a symbol called '§' that accounts for all special symbols. The total number of symbols in my alphabet is 96.

Originally I used a "for loop" with a "switch" statement to find all symbols and their occurrence. However, I found a problem reading special characters as new line etc.

Trying to optimize the code, I used the function *findstr* along with the function *length* to find the occurrence of each symbol. I had the same problem counting the special characters. Therefore as per Dr. He recommendation, all unknown symbols were assigned to a special character '§'. I implemented function *countocurrence(s,Nsym,sffile)*; to count all symbols. The occurrence and the relative frequency (we will use this number as the probability) is shown in the appropriate columns of Table 1.

I found that state.txt does not have all the 96 symbols from the alphabet created; therefore the occurrence, and hence, probability of some symbols is zero, leading the length to

"infinite" if we use the Shannon formula $l_i = \left[\log_2 \frac{1}{p_i} \right]$. One form of dealing with this

problem is to add at least one character to the training text. However I do not know what would be the optimization of doing this. I will investigate this at the end of my project.

order	symbol	Occur.	Rel. Freq	order	symbol	Occur.	Rel. Freq
			(prob)				(prob)
1	а	2817	0.0618	49	W	42	0.0009
2	b	374	0.0082	50	Х	1	2.2e-005
3	с	1121	0.0246	51	Y	19	0.0004
4	d	1364	0.0299	52	Z	0	0
5	e	4498	0.0987	53	Space	7561	0.1658
6	f	761	0.0167	54		519	0.0114
7	g	728	0.016	55	,	456	0.01
8	h	1575	0.0345	56	/	2	4.4e-005
9	i	2645	0.058	57	<	0	0
10	j	23	0.0005	58	>	0	0
11	k	180	0.0039	59	?	9	0.0002
12	1	1406	0.0308	60	;	10	0.0002
13	m	913	0.02	61	'	53	0.0012
14	n	2725	0.0598	62	:	37	0.0008
15	0	2788	0.0612	63		83	0.0018
16	р	709	0.0156	64	[61	0.0013
17	q	55	0.0012	65]	61	0.0013
18	r	2388	0.0524	66	{	0	0
19	s	2390	0.0524	67	}	0	0
20	t	3158	0.0693	68	\	0	0
21	u	943	0.0207	69		0	0
22	v	384	0.0084	70	!	4	0.0001
23	w	606	0.0133	71	@	0	0
24	Х	59	0.0013	72	#	0	0
25	у	515	0.0113	73	\$	2	4.4e-005
26	Z	33	0.0007	74	%	0	0
27	А	168	0.0037	75	^	0	0
28	В	28	0.0006	76	&	0	0
29	С	36	0.0008	77	*	2	4.4e-005
30	D	28	0.0006	78	(2	4.4e-005
31	E	17	0.0004	79)	2	4.4e-005
32	F	18	0.0004	80	-	31	0.0007
33	G	16	0.0004	81	=	0	0
34	Н	51	0.0011	82	_	0	0
35	I	139	0.003	83	+	0	0
36	J	7	0.0002	84	`	0	0
37	K	11	0.0002	85	~	0	0
38	L	9	0.0002	86	0	72	0.0016
39	М	27	0.0006	87	1	35	0.0008
40	Ν	44	0.001	88	2	16	0.0004
41	0	27	0.0006	89	3	7	0.0002
42	Р	39	0.0009	90	4	7	0.0002
43	Q	10	0.0002	91	5	11	0.0002
44	R	15	0.0003	92	6	6	0.0001
45	S	108	0.0024	93	7	2	4.4e-005
46	Т	101	0.0022	94	8	4	0.0001
47	U	35	0.0008	95	9	17	0.0004
48	V	3	0.0001	96	Schar	363	0.008

Table 1: Symbol Table, with occurrences and relative frequency

With this table, and the function in Appendix 1, **step 1** of the project is concluded.

Step 2 and 3: (Huffman Algorithm – average length)

The proposed Huffman algorithm works as follows: from top to bottom we search for the two minimum symbols with less probability. We mark those symbols as counted and create a new super symbol with both symbols and probability the sum of the

probabilities. Then repeat the process until completion of the unmarked symbols. If there are 5 symbols we have to do 4 comparisons.

Example:

symbols	prob
a	.25
b	.25
c	.20
d	.15
e	.15

Table 2: example Symbol Table, with probabilities

The algorithm finds **d** and then finds **e**. Assigns 0 to d and 1 to e

Then, update the symbol table with a super symbol with probability the sum of the probabilities of the found symbols. In order to make the algorithm work, I set the probabilities of the found symbols an impossible number (like 2) in order to be larger than the maximum probability (this is a way of marking them as counted).

prob	Code
.25	
.25	
.20	
2	0
2	1
.30	
	prob .25 .25 .20 2 .30

Table 3: example Symbol Table, with probabilities and code

Now it repeats itself. At the second round we find **c** and **a**. Assigns 0 to c and 1 to a

symbols	prob	Code
а	2	1
b	.25	
с	2	0
d	2	0
e	2	1
de	.30	
ca	.45	

Table 4: example Symbol Table, with probabilities and code

At the third round it finds **b** and **de**. Assigns 0 to b and 1 to both d and e.

symbols	prob	Code
a	2	1
b	2	0
с	2	0
d	2	10
e	2	11
de	2	
ca	.45	
bde	.55	

 Table 5: example Symbol Table, with probabilities and code

At the last round (N-1), it finds ca and bde. Assigns 0 to c and a and 1 to b, d and e

symbols	prob	Code
а	2	01
b	2	10
с	2	00
d	2	110
e	2	111
de	2	
ca	2	
bde	2	
cabde	1	

 Table 6: example Symbol Table, with probabilities and code

The final code is shown next:

symbols	code
a	01
b	10
с	00
d	110
e	111

Table 7: Huffman code for this example.

I coded this algorithm, and came out with the Huffman code. The Matlab code it is written in Appendix 1 (look for fhuffman function). The Huffman code for this project is written in Table 8. The hexadecimal representation of each codeword is displayed in Appendix 1.

<u>Observations</u>: Analyzing the code, I found that the maximum length is 21. It seems is a pretty big number for just 96 symbols. Let's find the average length or bit rate, defined as the summation of the product of the probability of a symbol with the length of the code for that symbol.

$$L = \sum_{i=1}^{N} p_i l_i = 4.4563 \text{ (Step 3)}$$

This seems ok with me, since if we look in Table 8, the symbols with more occurrences (space, e) have code length of 3. Other frequent symbols (a, i, n, o etc) have length of 4. The symbols with less frequency (c, f) have lengths of 5 and 6. Finally, symbols that do not appear (frequency 0) have lengths of 20 and 21.

I added to the training file the missing characters (Z, <, >, {, }, \, |, @, #, ^, &, =, _, +, `, ~), to see if we improved the efficiency and I found that in this case even when the maximum length reduces from 21 to 16; the bit rate or average length raises from 4.4563 to 4.4607. Therefore I think there is not optimization in adding the missing symbols to the training file.

order	symbol	Occur.	Huffman Code.	length		order	symbol	Occur.	Huffman Code.	length
1	a	2817	'1001'	4		49	W	42	'0110001011'	10
2	b	374	'1011000'	7		50	Х	1	'1101110101000001'	16
3	с	1121	'00100'	5		51	Y	19	'01100010001'	11
4	d	1364	'01101'	5		52	Z	0	'110111010100000011110'	21
5	e	4498	'000'	3		53	Space	7561	'111'	3
6	f	761	'101101'	6		54		519	'001010'	6
7	g	728	'101011'	6		55	,	456	'1101001'	7
8	h	1575	'10111'	5		56	/	2	'110111010100001'	15
9	i	2645	'0101'	4		57	<	0	'1101110101000000111111'	21
10	j	23	'10101000110'	11		58	>	0	'1101110101000000000'	20
11	k	180	'10101001'	8		59	?	9	'011000100001'	12
12	1	1406	'10100'	5		60	;	10	'011000101010'	12
13	m	913	'110101'	6		61	'	53	'1101000011'	10
14	n	2725	'0111'	4		62	:	37	'0110000111'	10
15	0	2788	'1000'	4		63	"	83	'101010000'	9
16	р	709	'011001'	6		64	[61	'1101110011'	10
17	q	55	'1101000111'	10		65		61	'1101110100'	10
18	r	2388	'0011'	4		66	{	0	'11011101010000000001'	20
19	S	2390	'0100'	4		67	}	0	1101110101000000010	20
20	t	3158	'1100'	4		68	\	0	1101110101000000011	20
21	u	943	'110110'	6		69		0	11011101010000000100	20
22	V	384	1011001	1		70	!	4		14
23	w	606	001011	0		/1	@	0		20
24	X	59	1101110001	10		72	# ¢	0	110111010000000110	20
25	y	215	1101111	/		75	\$	2		15
20	Z	169	101100011	0		74	%	0	<u>'11011101010000001111</u>	20
27	A	108	'11011100000'	0		75	P_	0	<u>'1101110101000001000</u>	20
20	D C	26	'01100000	10		70	<u>«</u>	0	'11011101000001001	20
29		28	'1101110000110	10		78	*	2	'01100010100000'	13
30	E D	17	'11011100001	11		70	(2	'011000101000001'	14
32	E F	17	'110111011111'	12		80)	31	'11011100101'	14
33	G	16	'110111010101'	12		81	_	0	'110111010000001010'	20
34	Н	51	'1101000010'	10		82		0	'11011101010000001011'	20
35	I	139	'01100000'	8		83	+	0	'11011101010000001100'	20
36	J	7	'1010100011111'	13		84		0	'11011101010000001101'	20
37	K	11	'011000101011'	12		85	~	0	'11011101010000001110'	20
38	L	9	'011000100000'	12		86	0	72	'011000010'	9
39	М	27	'11010001100'	11		87	1	35	'11011101110'	11
40	N	44	'1010100010'	10		88	2	16	'110111011000'	12
41	0	27	'11010001101'	11		89	3	7	'1101110010000'	13
42	Р	39	'0110001001'	10		90	4	7	'1101110010001'	13
43	0	10	'011000101001'	12		91	5	11	'101010001110'	12
44	R	15	'110111001001'	12		92	6	6	'1010100011110'	13
45	S	108	'110100010'	9		93	7	2	'01100010100010'	14
46	Т	101	'110100000'	9		94	8	4	'11011101010010'	14
47	U	35	'11011101101'	11	l	95	9	17	'110111011110'	12
48	V	3	'01100010100011'	14		96	Schar	363	'1010101'	7

 Table 8: Huffman code for the project showing the codeword and the length of each codeword.

APPENDIX 1

This is the Matlab program I wrote for this project. It consists of several functions I will summarize here: main() is the main function, it calls the other functions and perform the display and final computations. readfile() this function read the characters of the training text file enumsymbols() this function creates the dictionary of symbols countocurrence(s,Nsym,sffile) this function counts the symbols from the training file and gives how many symbols of the dictionary are in the file.

fhuffman(st,p) this function performs the Huffman algorithm

%mreadf05 (The last number is the version of this code) %read all characters in a file, perform the Huffman code and display %Created 2/3/2005 by Luis M Vicente SN: 945-995 %Deficiencies in version 04: I am going to do a new version because I want %to read the special characters in a supper symbol. I can not use n(i) = %length(findstr(sffile,s(i))); because it does not account for unknown %symbols. I will use a switch statement or something else and all unnkwon %characters will be stored in a special character symbol %Creating version mreadf05
%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%
%Function Read file and store it in array sffile. N is the # of characters [sffile,N] = readfile();
%Function to create a look up table of all symbols [s,Nsym]=enumsymbols();
%Function to count the occurrence of each symbol in the training set %store it in a vector that is sincronized with the symbol vector n = countocurrence(s,Nsym,sffile);
nt=n'; %column vector st=s'; %column vector
%Find the probabilities $p = nt/N;$
%This is just to display the symbol table, the occurrence, and the %probability mcell=cell(length(st),2); mcell(:,1)=cellstr(st); mcell(:,2)=num2cell(nt); mcell(:,3)=num2cell(p);
%Display symbol, occurrence and probability disp(' '); disp(' Symbol, [Occurrence], [Probability]'); mcell
disp('End of Step 1 of the project. '); disp('To continue with Step 2, please press "Enter"'); pause
%Step 2 Function to create Huffmann code

[codem] = fhuffman(st,p);
%Adding the code to display later mcell(:,4)=codem(:,2);
%sum of all probabilities, =1?? disp(' '); disp(['Sum of all probabilities : ', num2str(sum(p,1))]);
%computing the length of each symbol for i=1:Nsym len(i,1) = length(char(codem(i,2))); end
%Display symbol, occurrence, probability, codeword and length disp(' '); disp(' Symbol, [Occurrence], [Probability] [codeword] [hex] [length]');
mcell(:,5)=cellstr(dec2hex(bin2dec(char(mcell(:,4))))); mcell(:,6)=num2cell(len)
%Bit rate or Average Length Lengthaverage = sum(p.*len ,1); disp(['Bit rate or Average Length : ', num2str(Lengthaverage)]);
%Max length maxlength = max(len); disp(['Max codeword Length : ', num2str(maxlength)]);
%Lets find how many words have length of 16 15 14 etc. for i=1:maxlength maxl(i,1)=0; for j=1:Nsym if len(j)==i maxl(i,1)=maxl(i,1)+1; end end end
%Take out the semicolon if you want to display how many words have lengths %of 1 2 3 up to max length maxl;
%This is a checkup to see if the characters read are the same as the %total characters in the training file if sum(n,2) ~= N disp('Error counting characters in this file'); else disp(['Characters read: ', num2str(N)]); end
disp(['END OF MAIN %%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%
%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%
%Filename strf = 'state.txt';
%Open the file fid = fopen(strf); if(fid ~= -1) disp(['File ',strf,' succesfully opened']); else disp(['ERROR OPENING FILE ',strf]); return

end

%Read text from file and count characters [sffile,N] = fscanf(fid,'%c'); %closing the file if(fclose(fid) ==0) disp(['File ',strf,' succesfully closed']); else disp(['ERROR CLOSING File ',strf]); end function [s,Nsym]=enumsymbols() %Enumeration of symbols/ Creation of the alphabet %OUTPUT: s (is a cell of strings) and Nsym (integer) the number of %symbols in the alphabet %to save space I wrote them in the way it is shown. %Orthodox programmers forgive me!! $s(1) = {'a'}; s(2) = {'b'}; s(3) = {'c'}; s(4) = {'d'};$ $s(5) = { 'e' }; s(6) = { 'f' }; s(7) = { 'g' }; s(8) = { 'h' };$ $s(9) = {'i'}; s(10) = {'j'}; s(11) = {'k'}; s(12) = {'l'};$ $s(13) = {'m'}; s(14) = {'n'}; s(15) = {'o'}; s(16) = {'p'};$ $s(17) = \{ 'q' \}; s(18) = \{ 'r' \}; s(19) = \{ 's' \}; s(20) = \{ 't' \};$ $s(21) = {'u'}; s(22) = {'v'}; s(23) = {'w'}; s(24) = {'x'};$ s(25) = {'y'};s(26) = {'z'};s(27) = {'A'};s(28) = {'B'}; s(29) = {'C'};s(30) = {'D'};s(31) = {'E'};s(32) = {'F'}; $s(33) = {'G'}; s(34) = {'H'}; s(35) = {'I'}; s(36) = {'J'};$ $s(37) = {'K'}; s(38) = {'L'}; s(39) = {'M'}; s(40) = {'N'};$ $s(41) = \{'O'\}; s(42) = \{'P'\}; s(43) = \{'Q'\}; s(44) = \{'R'\}; \\ s(45) = \{'S'\}; s(46) = \{'T'\}; s(47) = \{'U'\}; s(48) = \{'V'\};$ $s(49) = {'W'}; s(50) = {'X'}; s(51) = {'Y'}; s(52) = {'Z'};$ s(53) = {' '}; %Space character $s(54) = \{'.'\}; s(55) = \{','\}; s(56) = \{'/'\}; s(57) = \{'<'\}; s(58) = \{'>'\}; s(59) = \{'?'\}; s(60) = \{';'\}; s(61) = \{''''\};$ $s(62) = {':'}; s(63) = {'''}; s(64) = {'[']; s(65) = {']'};$ $s(66) = \{ \{ \}; s(67) = \{ \} \}; s(68) = \{ \{ \}; s(69) = \{ \{ \} \}; s(69) = \{ \{ \} \}; s(69) = \{ \{ \{ \} \}$ $s(70) = {'!'}; s(71) = {'@'}; s(72) = {'#'}; s(73) = {'$'};$ $s(74) = {'\%'}; s(75) = {'^{+}}; s(76) = {'\&'}; s(77) = {'*'};$ $s(78) = {(')}; s(79) = {()}; s(80) = {(-)}; s(81) = {(=)};$ $s(82) = {'-}; s(83) = {'+}; s(84) = {''}; s(85) = {'-'};$ $s(86) = {0}; s(87) = {1}; s(88) = {2}; s(89) = {3};$ $s(90) = {'4'}; s(91) = {'5'}; s(92) = {'6'}; s(93) = {'7'};$ $s(94) = \{8\}; s(95) = \{9\};$ $s(96) = {'\S'};$ %This accounts for all special characters %For the near future, to add new characters to the alphabet % s(96) = (char(10)); %New Line? % s(97) = (char(13)); %New Line? % s(98) = (char(151)); %Strange square Nsym = length(s); function [n] = countocurrence(s,Nsym,sffile) %function [n] = countocurrence(s,Nsym,sffile) %Count the occurrence of each symbol in the training set except the special %character symbol %store it in a vector that is sincronized with the symbol vector %INPUTS: s (cell of characters - symbols), Nsym (integer - number of %symbols), sifile (char array - text having the symbols to count) %OUTPUT: n (array of integers -Count of each symbol indexed) for i=1:Nsym-1 %Count the known symbols n(i) = length(findstr(sffile,char(s(i)))); end %Unknown symbols = total symbols - known symbols n(Nsym)=length(sffile)-sum(n,2);

```
function [codem] = fhuffman(s,p)
%function [codem] = fhuffman(s,p)
%This algoritm creates a Huffman code from s (a cell of strings) and a
%p (vector of floats) which is synchronized with s. meaning p(1) is the
%probability of symbol s(1).
%This algorithm uses string comparison, therefore the condition for vector
%s is that all characters should be SINGLE CHARACTER and different from
%each other.
%The outpus is a cell with two columns, first has the symbols and the
%second column has the codeword
%Example fhuffman({'a';'b';'c';'d';'e'},[.25;.25;.2;.15;.15])
%length of the diccionary
len = length(s);
%CodeMatrix. First column is initialized to the symbols in s,
%second column is emptied will be used to store the Huffman code as a
%binary string (string of 0's and 1's)
codem=cell(len,2);
codem(:,1)=cellstr(s);
codem(:,2)={"};
%We will modify the p vector, so we create a copy to work with it
pwork = p;
%We will modify the s vector, so we create a copy to work with it
swork = s;
%Find first two symbols with the smallest probability
%take that value from the vector, and repeat
%until we have counted all probabilities.
for i=1:len-1
  %first minimum symbol
  [pmin1,ind0]=min(pwork);
  pwork(ind0)=2; %Replace with an impossible probability
  %second minimum symbol
  [pmin2,ind1]=min(pwork);
  pwork(ind1)=2; %Replace with an impossible probability
  %Add probabilities
  pnew = pmin1+pmin2;
  %update pwork adding at the end of the vector the new
  %supersymbol and the new probability
  lsw = length(swork);
  pwork(lsw+1,1)=pnew;
  %Assign bits to the selected symbols.
  for j=1:len
    if double(findstr(char(swork(ind0)),char(s(j)))) >=1
       %disp(['Assign 0 to ',s(j)])
       codem(j,2)={['0',char(codem(j,2))]};
    end
    if double(findstr(char(swork(ind1)),char(s(j)))) >=1
       %disp(['Assign 1 to ',s(j)])
       codem(j,2)={['1',char(codem(j,2))]};
    end
  end
  %Create new super symbol
  ssym = {[char(swork(ind0)),char(swork(ind1))]};
  swork(lsw+1,1)=ssym;
end
%EOM
```

The actual screenshot when running this program in Matlab is shown next:

START OF MA	IN %%%%%%%%	***************************************
File state.	txt succes	fully opened
File state.	txt succes	fully closed
Crmbol	[Ogguerong	ol [Drobobility]
Symbol,	loccurrence	e], [Probability]
mcell =		
'a'	[2817]	[0.0618]
'b'	[374]	[0.0082]
'c'	[1121]	[0.0246]
'd'	[1364]	
'e'	[4498]	
' 1 '	[701] [728]	
9 'h'	[1575]	[0.0345]
'i'	[2645]	0.05801
'j'	[23]	[5.0447e-004]
' k '	[180]	[0.0039]
'1'	[1406]	[0.0308]
'm'	[913]	[0.0200]
'n'	[2725]	[0.0598]
'0'	[2788]	
'P'	[709]	
'q'	[55]	
 'e'	[2300]	
'+'	[3158]	[0.0693]
'u'	[943]	0.02071
'v'	[384]	[0.0084]
'w'	[606]	[0.0133]
'x'	[59]	[0.0013]
'Y'	[515]	[0.0113]
'z'	[33]	[7.2381e-004]
'A'	[168]	[0.0037]
'B'	[28]	[6.1414e-004]
101	[36]	[7.8961e-004]
ע ידי	[20]	$[3, 7287_{-004}]$
<u>ה</u> יקי	[18]	[3, 9481e-0.04]
'G '	[16]	[3.5094e-004]
'H'	[51]	[0.0011]
'I'	[139]	[0.0030]
'J'	[7]	[1.5354e-004]
'K '	[11]	[2.4127e-004]
'L'	[9]	[1.9740e-004]
' M '	[27]	[5.9221e-004]
'N'	[44]	[9.6508e-004]
יסי יסי	[27]	[5.9221e-004]
'0'	[10]	[2.1934e-004]
'R'	[15]	[3.2901e-004]
'S'	[108]	[0.0024]
'T'	[101]	[0.0022]
יטי	[35]	[7.6768e-004]
' V '	[3]	[6.5801e-005]
' W '	[42]	[9.2121e-004]
'X'		[2.1934e-005]
· Y ·	[17]	[4.10/4e-UU4]
· Z ·	L U] [7561]	
	[519]	[0.114]
	[456]	0.0100]
'/'	[2]	[4.3867e-005]
' < '	[0]	[0]
'>'	[0]	[0]

'?'	[9]	[1.9740e-004]			
1 ; 1	[10]	[2 1934e - 0.04]			
	[55]				
':'	[37]	[8.1155e-004]			
	[83]	[0.0018]			
'['	[61]	[0.0013]			
1	[61]	[0.0013]			
1	[0]	[0:0013]			
1		[0]			
'}'	[0]	[0]			
'\'	[0]	[0]			
1 1	[0]	[0]			
	[4]	[8 7735e-005]			
!@!	[0]	[0]			
		[0]			
'#'		[0]			
'\$'	[2]	[4.3867e-005]			
181	[0]	[0]			
1 ^ 1	[0]	[0]			
1.5.1	[0]	[0]			
1*1	[2]				
		[4.38678-005]			
'('	[2]	[4.3867e-005]			
')'	[2]	[4.3867e-005]			
' = '	[31]	[6.7994e-004]			
' = '	[0]	[0]			
	[0]				
-					
'+'					
1 1	[0]	[0]			
'~'	[0]	[0]			
'0'	[72]	[0.0016]			
11	[35]	[7 6768e-004]			
1.21	[16]				
2	[10]	[3.50940-004]			
'3'	[7]	[1.5354e-004]			
' 4 '	[7]	[1.5354e-004]			
'5'	[11]	[2.4127e-004]			
'6'	[6]	[1.3160e-004]			
171	[2]	[4 38670-005]			
101					
.8.	[4]	[8.7/350-005]			
'9'	[17]	[3.7287e-004]			
'§'	[363]	[0.0080]			
End of St	ep 1 of th	e project.			
To cor	ntinue with	Step 2 please p	regg "Enter"		
	iciliac wich	beep 2, picase p	ICBS ENCCI		
G	1				
Sum of al	ll probabil	ities : 1			
Symbol	L, [Occurre	nce], [Probabilit	y] [codeword]	[hex]	[length]
-	, -				- 5 -
mcell -					
MCEIT =					
'a'	[2817]	L 0.0618]	'1001'	'000009'	[4]
'b'	[374]	[0.0082]	'1011000'	'000058'	[7]
'c'	[1121]	[0.02461	'00100'	'000004'	[5]
יהי	[1364]	[0.0299]	'01101'	יתחחחחי	[5]
lal	[1100]		10001	1000001	[2]
	[7/1]		101101		
, Γ,	[/61]	[0.0167]	, TOTIOT,	·00002D·	[6]
'g'	[728]	L 0.0160]	'101011'	'00002B'	[6]
'h'	[1575]	[0.0345]	'10111'	'000017'	[5]
'i'	[2645]	[0.0580]	'0101'	'000005'	[4]
141	[23]	[5.0447e - 0.041]	10101000110	000546	[11]
ן ו - 1 ו	[100]		101010011	10000701	[9]
K.			101001	10000A9	
' ['] '	[1406]	[0.0308]	. TOTOO .	'000014'	[5]
'm'	[913]	L 0.0200]	'110101'	'000035'	[6]
'n'	[2725]	[0.0598]	'0111'	'000007'	[4]
'o'	[2788]	[0.0612]	'1000'	'000008'	[4]
'n'	[709]	[0.0156]	011001	10000101	[6]
P	[[[[]]		11010001111	10002	[10]
. Ч	[22]		TIDIOODITI	000347	
'r'	[2388]	L 0.0524]	,00TT ,	'000003'	[4]
's'	[2390]	[0.0524]	'0100'	'000004'	[4]
't'	[3158]	[0.06931	'1100'	'00000C'	[4]
'11'	[943]	[0.0207]	'110110'	0000361	[6]
I			1011001	10000501	[7]
1 7 7	[224]	0 0084	' ' ' ' ' ' ' ' ' ' ' ' ' ' ' ' ' ' ' '		
'V'	[384]		'1011001'	000059	

'x'	[59]	[0.0013]	'1101110001'	'000371'	[10]
'v'	[515]	[0.0113]	'1101111'	'00006F'	[7]
2	[33]	$\begin{bmatrix} 7 & 23816 - 0041 \end{bmatrix}$	11011101011	10006FB1	[11]
17.1	[169]		101100011	10000621	[0]
A	[100]	$\begin{bmatrix} 0.0037 \end{bmatrix}$	11011100001	10006001	
В	[20]	[8.1414e-004]		10000E0	
101	[30]	[7.89610-004]	. 0110000110	000186	
· D ·		[6.1414e-004]		.0006E1.	
'E'	[17]	[3.7287e-004]	'110111011001'	'000DD9'	[12]
' F' '	[18]	[3.9481e-004]	'1101110111111'	'000DDF'	[12]
'G'	[16]	[3.5094e-004]	'110111010101'	'000DD5'	[12]
'H'	[51]	[0.0011]	'1101000010'	'000342'	[10]
'I'	[139]	[0.0030]	'01100000'	'000060'	[8]
'J'	[7]	[1.5354e-004]	'1010100011111'	'00151F'	[13]
'K '	[11]	[2.4127e-004]	'011000101011'	'00062B'	[12]
'L'	[9]	[1.9740e-004]	'011000100000'	'000620'	[12]
' M '	[27]	[5.9221e-004]	'11010001100'	'00068C'	[11]
' N '	[44]	[9.6508e-004]	'1010100010'	'0002A2'	[10]
'0'	[27]	[5.9221e-004]	'11010001101'	'00068D'	[11]
'P'	[39]	[8.5541e-004]	'0110001001'	'000189'	[10]
'0'	[10]	[2.1934e-004]	'011000101001'	'000629'	[12]
'R'	[15]	[3.2901e-004]	110111001001	'000DC9'	[12]
151	[108]	[0.0024]	110100010'	'0001A2'	[9]
יידי	[101]		110100000'	000120	[9]
'11'	[35]	$\begin{bmatrix} 7 & 6768e - 0.04 \end{bmatrix}$	11011101101	'0006FD'	[11]
1771	[3]			1001050	
V 1 107 1	[42]	[0.3101 - 004]		10001000	
· W ·		[9.21210-004]		.00018B	
· X ·				·00DD41·	[10]
Ϋ́Υ,	[19]	[4.16/4e-004]	.01100010001,	'000311'	
'Z'		[0]	[1x21 char]	'IBA81E'	[21]
	[7561]	[0.1658]	'111'	000007	[3]
'.'	[519]	[0.0114]	'001010'	'00000A'	[6]
', '	[456]	[0.0100]	'1101001'	'000069'	[7]
'/'	[2]	[4.3867e-005]	'110111010100001'	'006EA1'	[15]
' < '	[0]	[0]	[1x21 char]	'1BA81F'	[21]
' > '	[0]	[0]	[1x20 char]	'0DD400'	[20]
'?'	[9]	[1.9740e-004]	'011000100001'	'000621'	[12]
';'	[10]	[2.1934e-004]	'011000101010'	'00062A'	[12]
	[53]	[0.0012]	'1101000011'	'000343'	[10]
':'	[37]	[8.1155e-004]	'0110000111'	'000187'	[10]
	[83]	[0.0018]	101010000'	000150'	[9]
יןי	[61]	[0.0013]	'1101110011'	000373	[10]
יןי	[61]	[0.0013]	1101110100'	000374	[10]
, , ,	[0]	[0]	[1x20 char]	' 00000/1	[20]
ן י ג י	[0]	[0]	[1x20 char]	י 402 מסטי	[20]
ر ب\ب		[0]	[1x20 char]	100402	[20]
		[0]	[1x20 char]	100400	[20]
				10027511	[20]
		[8.//350-005]		003/51	[14]
		[0]	[1x20 Char]	· 0DD405 ·	[20]
· # ·			[IX20 Char]	000406	[20]
· Ş '		[4.3007e-005]	TINTITUTUTUTUTU	UUBEA6'	[15]
			[1x20 char]	' UDD40'7 '	
			[1x20 char]	'UDD408'	[20]
`&_'			[1x20 char]	'UDD409'	
1 * 1	L 2]	[4.3867e-005]	. TIOTITOTOTOOJJJ ,	'UU6EA7'	[15]
'('	[2]	[4.3867e-005]	.01100010100000.	'0018A0'	
')'	L 2]	[4.3867e-005]	'01100010100001'	'0018A1'	[14]
' - '	[31]	[6.7994e-004]	'11011100101'	'0006E5'	[11]
' = '	[0]	[0]	[1x20 char]	'0DD40A'	[20]
'_'	[0]	[0]	[1x20 char]	'0DD40B'	[20]
'+'	[0]	[0]	[1x20 char]	'0DD40C'	[20]
1 2 1	[0]	[0]	[1x20 char]	'0DD40D'	[20]
'~'	[0]	[0]	[1x20 char]	'0DD40E'	[20]
'0'	[72]	[0.0016]	'011000010'	'0000C2'	[9]
'1'	[35]	[7.6768e-004]	'11011101110'	'0006EE'	[11]
'2'	[16]	[3.5094e-004]	'110111011000'	'000DD8'	[12]
'3'	[7]	[1.5354e-004]	'1101110010000'	'001B90'	[13]
'4'	[7]	[1.5354e-004]	'1101110010001'	'001B91'	[13]
'5'	[11]	[2.4127e-004]	'101010001110'	'000A8E'	[12]
'6'	[6]	[1.3160e-004]	'1010100011110'	'00151E'	[13]
'7'	[2]	[4.3867e-005]	'01100010100010'	'0018A2'	[14]
' 8 '	[4]	[8.7735e-005]	'11011101010010'	'003752'	[14]

	'9'	[17]	[3.7287e-004]		'110111011110'	'000dde'	[12]		
	'§'	[363]	[0.0080]	'1010101'	'000055'	[7]		
Bit rate or Average Length : 4.4563 Max codeword Length : 21 Characters read: 45592										
END OF MAIN %%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%										
End Of Project										